## **Randomized Algorithms for Geometric Structure**

- 1. Analyze the expected time complexity of the randomized Quick-Sort and explain the relation between the sorting and geometric partition.
- 2. Define the conflict list for an interval and describe the randomized incremental version of Quick-Sort.
- 3. Analyze the expected time complexity of the randomized incremental version of Quick-Sort (Backward Analysis)
- 4. Use a randomized binary tree to develop an on-line version of quick-sort and analyze the expected time complexity
- 5. Prove that given a set N of n line segments with total k intersections and an *i*-element subset  $N^i$  of N, the expected number of trapezoids in the vertical trapezoidal decomposition  $H(N^i)$  of  $N^i$  is  $O(i + ki^2/n^2)$ .
- 6. Define conflict relations between a newly inserted segment and the current trapezoidal decomposition  $H(N^i)$ , and describe how to insert a new segment.
- 7. Analyze the expected time of inserting a line segment into  $H(N^i)$ , and the total expected time for constructing the vertical trapezoidal decomposition.
- 8. Describe how to use a history graph to develop an on-line algorithm for the vertical trapezoidal decomposition and analyze the expected time complexity.
- 9. Please compare conflict graphs and history graphs.
- 10. Regarding the paper "Kenneth L. Clarkson, Kurt Mehlhorn, and Raimund Seidel Four, Results on Randomized Incremental Construction," define a configuration, conflict relations, and history, and give one example, e.g., vertical trapezoidal decomposition.
- 11. Please analyze the expected number of conflict relation between the configurations in  $H_r$  (the history for the first r element) and  $x_r$  (the  $r^{\text{th}}$  element
- 12. Please analyze the size of conflict history.

## Abstract Voronoi diagram

- 1. Define abstract Voronoi diagrams, describe the motivation, and list several examples.
- 2. Let  $(S, \mathcal{J})$  be a bisecting curve system. Please prove that the following assertions are equivalent.
  - If p, q, and r are pairwise different sites in S, then  $D(p,q) \cap D(q,r) \subseteq D(p,r)$  (Transitivity)
  - For each nonempty subset  $S' \subseteq S$ ,  $R^2 = \bigcup_{p \in s'} \overline{\operatorname{VR}(p, S')}$
- 3. Please argue that for checking an admissible system, it is enough to check all subset of 3 sites.
- 4. Define a conflict graph for the incremental construction of AVD, and prove that local test is enough, i.e.,  $e \cap \operatorname{VR}(t, R \cup \{t\}) = e \cap \operatorname{R}(t, \{p, q, r\})$ , where  $R \subseteq S, t \in S \setminus R$ , and e is the Voronoi edge between  $\operatorname{VR}(p, R)$ and  $\operatorname{VR}(q, R)$ .
- 5. Describe how to compute  $V(R \cup \{s\})$  from V(R)
- 6. Describe how to update the conflict graph, i.e., computing  $G(R \cup \{s\})$  from G(R).
- 7. Prove that the time to compute  $V(R \cup \{s\})$  and  $G(R \cup \{s\})$  from V(R) and G(R) is linear to the number of structural changes.
- 8. Use the general ideas of randomized geometric algorithms to analyze the expected time complexity for the randomized incremental construction of AVD.