# Discrete and Computational Geometry, SS 14 <br> Exercise Sheet " 4 ": Randomized Algorithms for Geometric Structures II <br> University of Bonn, Department of Computer Science I 

- Written solutions have to be prepared until Tuesday 13th of May, 14:00 pm. There will be a letterbox in the LBH building.
- You may work in groups of at most two participants.
- Please contact Hilko Delonge, hilko.delonge@uni-bonn.de, if you want to participate and have not yet signed up for one of the exercise groups.
- If you are not yet subscribed to the mailing list, please do so at https://lists.iai.uni-bonn.de/mailman/listinfo.cgi/lc-dcgeom


## Exercise 10: Planar Convex Hull by Conflict Lists (4 Points)

Given a set $N$ of $n$ points in the plane, a convex hull $H(N)$ of $N$ is a minimal convex polygon containing $N$, Let $S_{1}, S_{2}, \ldots, S_{n}$ be a random sequence of $N$, and let $N^{i}$ be $\left\{S_{1}, S_{2}, \ldots, S_{i}\right\}$. Please develop a randomized algorithm to construct $H(N)$ by computing $H\left(N^{3}\right), H\left(N^{4}\right), \ldots, H\left(N^{n}\right)$ iteratively using the conflict lists. In other words, for $i \geq 3$, obtain $H\left(N^{i+1}\right)$ from $H\left(N^{i}\right)$ by adding $S_{i+1}$.

1. Describe the insertion of $S_{i+1}$
2. Define a conflict relation between an edge of $H\left(N^{i}\right)$ and a point in $N \backslash N^{i}$
3. Prove the expected cost of inserting $S^{i+1}$ to be $O\left(\frac{n}{i+1}\right)$ and the expected cost of construction $H(N)$ to be $O(n \log n)$

## Exercise 11: Triangulation (History Graph)

Given a set $N$ of $n$ points in the plane, a triangulation $H(N)$ of $N$ is a maximal planar straight-light graph, i.e., every edge is a straight-line segment, and no edge can be added to main the planarity. Let $S_{1}, S_{2}, \ldots, S_{n}$ be a random sequence of $N$, and let $N^{i}$ be $\left\{S_{1}, S_{2}, \ldots, S_{i}\right\}$. Please develop a randomized algorithm to construct $H(N)$ by computing $H\left(N^{3}\right), H\left(N^{4}\right), \ldots, H\left(N^{n}\right)$ iteratively using the history graph. In other words, for $i \geq 3$, obtain $H\left(N^{i+1}\right)$ from $H\left(N^{i}\right)$ by adding $S_{i+1}$.

1. Describe the parent and child relation in the history graph.
2. Describe the insertion of $S_{i+1}$ using the history graph.
3. Prove the expected cost of inserting $S^{i+1}$ to be $O\left(\frac{n}{i+1}\right)$ and the expected cost of construction $T(N)$ to be $O(n \log n)$

## Exercise 12: Planar Convex Hull by History Graph (4 Points)

Given a set $N$ of $n$ points in the plane, a convex hull $H(N)$ of $N$ is a minimal convex polygon containing $N$, Let $S_{1}, S_{2}, \ldots, S_{n}$ be a random sequence of $N$, and let $N^{i}$ be $\left\{S_{1}, S_{2}, \ldots, S_{i}\right\}$. Please develop a randomized algorithm to construct $H(N)$ by computing $H\left(N^{3}\right), H\left(N^{4}\right), \ldots, H\left(N^{n}\right)$ iteratively using the history graph. In other words, for $i \geq 3$, obtain $H\left(N^{i+1}\right)$ from $H\left(N^{i}\right)$ by adding $S_{i+1}$.

1. Describe the parent and child relation in the history graph.
2. Describe the insertion of $S_{i+1}$ using the history graph.
3. Prove the expected cost of inserting $S^{i+1}$ to be $O\left(\frac{n}{i+1}\right)$ and the expected cost of construction $T(N)$ to be $O(n \log n)$
