

Recap:

WSPD w.r.t.  $s$

$(A_1, B_1), (A_2, B_2), \dots, (A_m, B_m)$

Theorem 19  $m \in O(s^d d^{\frac{d}{2}} n)$

comp. time  $O(d \cdot n \log n + s^d d^{\frac{d}{2}} \cdot n)$

$\nabla$ -Spanner  $P, q \in S$   $|N_G^q| \leq t(p, q)$

WSPD w.r.t.  $s \rightarrow$   
Spanner const.

$$t = \frac{5s^d}{s-4}$$

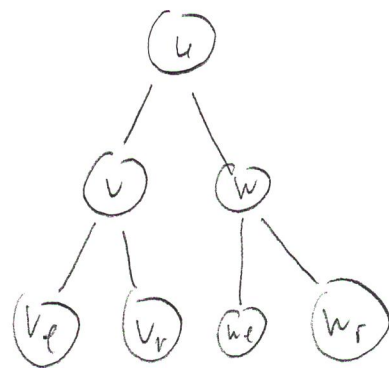
Constructing WSPD

(Partial) SPST Trees  $\rightarrow$  SPST Tree

Find Pairs Procedure

Find Pairs  $(v, w)$

$$L := \max \{ L_{\max}(R(s_v)), L_{\max}(R(s_w)) \}$$
$$r = \frac{\sqrt{d}}{2} L$$



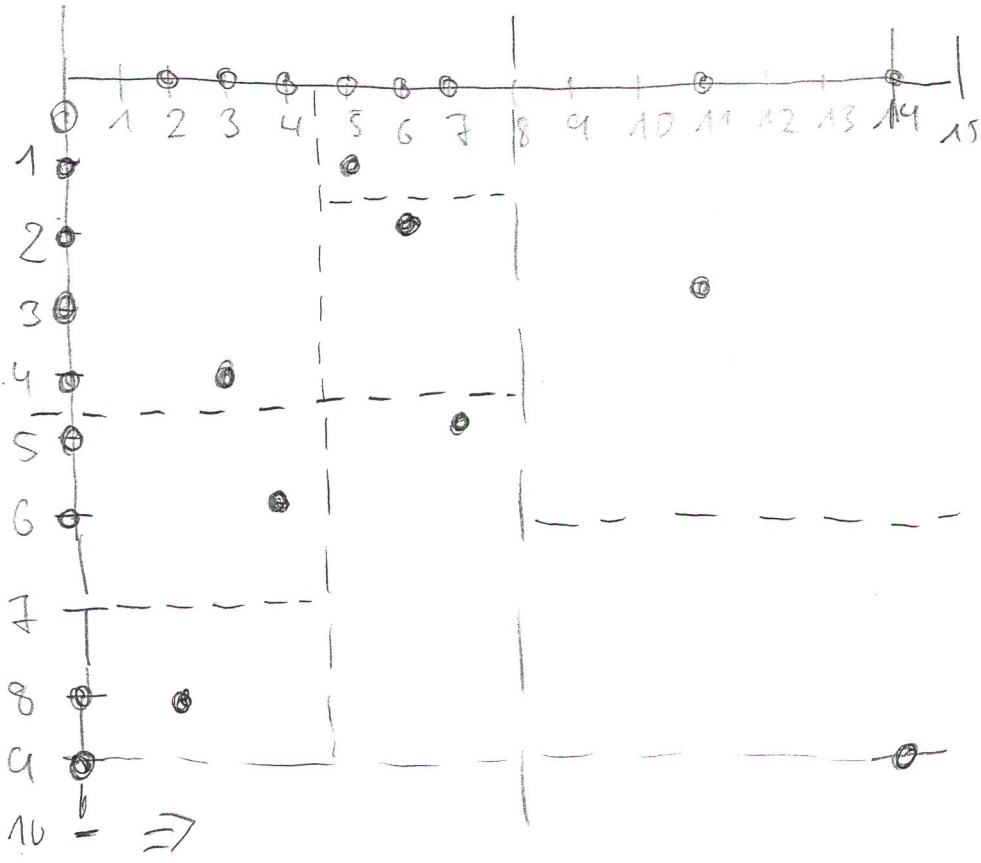
If  $(s_v, s_w)$  well-sep. w.r.t.  $s$  report  $(v, w)$

Else if  $L_{\max}(R(s_v)) > L_{\max}(R(s_w))$

Then Find Pairs  $(v_l, w)$ , Find Pairs  $(v_r, w)$

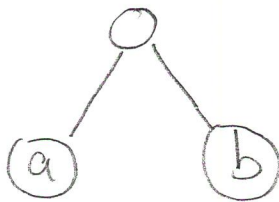
Else Find Pairs  $(v, w_l)$ , Find Pairs  $(v, w_r)$

# Partial Sperner Tree Construction



1.  $(2,8) (3,4) (4,6) (5,1) (6,2) (7,5) (11,3) (14,9)$   $S^1$
2.  $(5,1), (6,2) (11,3) (3,4) (7,5) (4,6) (2,8) (14,9)$   $S^2$

$L_{max} = 12$   
 $x = 8$

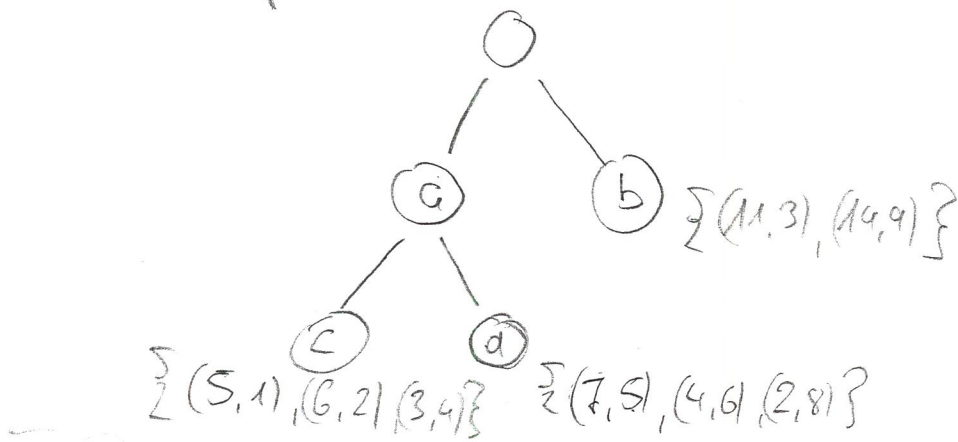


- Use copies of the bits
- Cross references between points!
- Doubly connected  $\Leftrightarrow (2,8) \leftrightarrow (6,2)$   
 $(7,5) \leftrightarrow (4,6)$
- $\{(11,3), (14,9)\}$

Step 1 Find points  $(11,3), (14,9)$  Hyperplane  $x=8$

- 2 Steps from right  $\leadsto$  One leaf!
- Delete points  $\{(11,3)\}$  and  $\{(14,9)\}$  out of copies for this phase!
- Cop's contains cross references

Step 2 for a



$\{ (2,8), (3,4), (4,6), (5,1), (6,7), (7,5) \}$

$\{ (5,1), (6,2), (3,4), (7,5), (4,6), (2,8) \}$

$\Rightarrow L_{max} = 7 \quad Y = 4,5$

Cross references to original list is still there ✓

3 Steps

Use original sets ✓

Leafs  $u$  and sets  $L_u$

Initial ~~ed~~ lists  $L_u^1, L_u^2$  for  $c, d, b = u$

For 1. (and 2.)

move through original lists and insert points in  $L_c^1, L_d^1, L_b^1$  at the end of the lists ✓

Example:

$L_c^1 = (3,4), (5,1), (6,2)$  and so on ✓

$L_d^1 = (2,8)$

$L_b^1 = (4,6)$

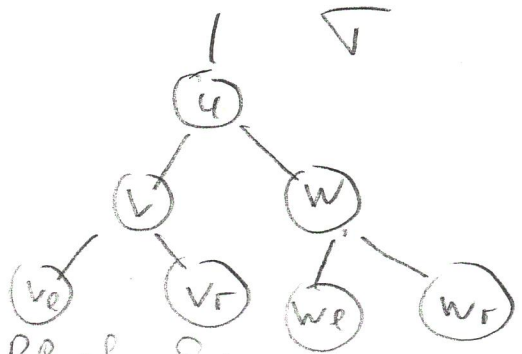
# Computing a NSPO of $S$ w.r.t. $s$

Given Split Tree  $\mathcal{T}$  of  $S$ ,  $s > 0$ ,  $s \in \mathbb{R}$

For each internal node  $u \in \mathcal{T}$  with children  $v, w$   
in node procedure

Find Pairs  $(v, w)$

$L_{\max}(R(S_v))$  longest diameter of BB of  $S_v$



(well-separated w.r.t. bounding boxes  $R(S_v) \in O(d)$   
w.r.t "radius"  $\frac{L_{\max}(R(S_v))}{2} \geq \frac{L_{\max}(R(S_w))}{2}$ )

Find Pairs  $(v, w)$

$$r = \frac{\sqrt{d}}{2} L \quad \left[ \begin{array}{|c|} \hline r \\ \hline \end{array} \right] \left[ \begin{array}{|c|} \hline r \\ \hline \end{array} \right] L$$

If  $S_v, S_w$  are well separated w.r.t  $s$  report  $(v, w)$

Else if  $L_{\max}(R(S_v)) \geq L_{\max}(R(S_w))$   $(S_v, S_w)$

Then Find Pairs  $(v_l, w)$ , Find Pairs  $(v_r, w)$

Else Find Pairs  $(v, w_l)$ , Find Pairs  $(v, w_r)$

Procedure terminates: Sets get smaller; for example

$$|S_{v_l}| \cdot |S_{w_l}| < |S_v| \cdot |S_w| \text{ and } |S_{v_r}| \cdot |S_{w_r}| < |S_v| \cdot |S_w|$$

$$L_{\max}(R(S_w)) \geq L_{\max}(R(S_v))$$

Lemma 20 The greedy procedure

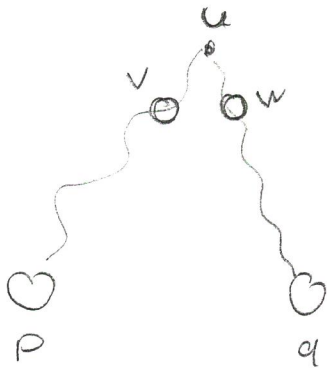
results in a WSPD w.r.t.  $S$

$$\{S_{v_1}, S_{w_1}\}, \{S_{v_2}, S_{w_2}\}, \dots, \{S_{v_n}, S_{w_n}\}$$

Proof: Vertices  $\rightarrow$  Leaves  $\times 2$ , size  $1$ , well-separated

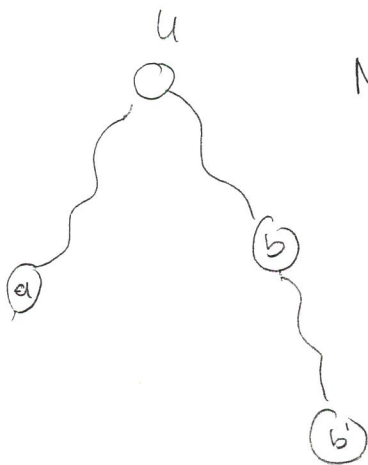
$p, q$  distinct points of  $S$

Exists  $\{S_{v_i}, S_{w_i}\}$  with  $p \in S_{v_i}, q \in S_{w_i}$ ?



Backwards from the leaves, lowest common ancestor  $u$  of  $p$  and  $q$  exists. Guarantees existence  $v$ . (lca)

Uniqueness:



More reports:

$$(a, b) \text{ and } (a, b') \Downarrow$$

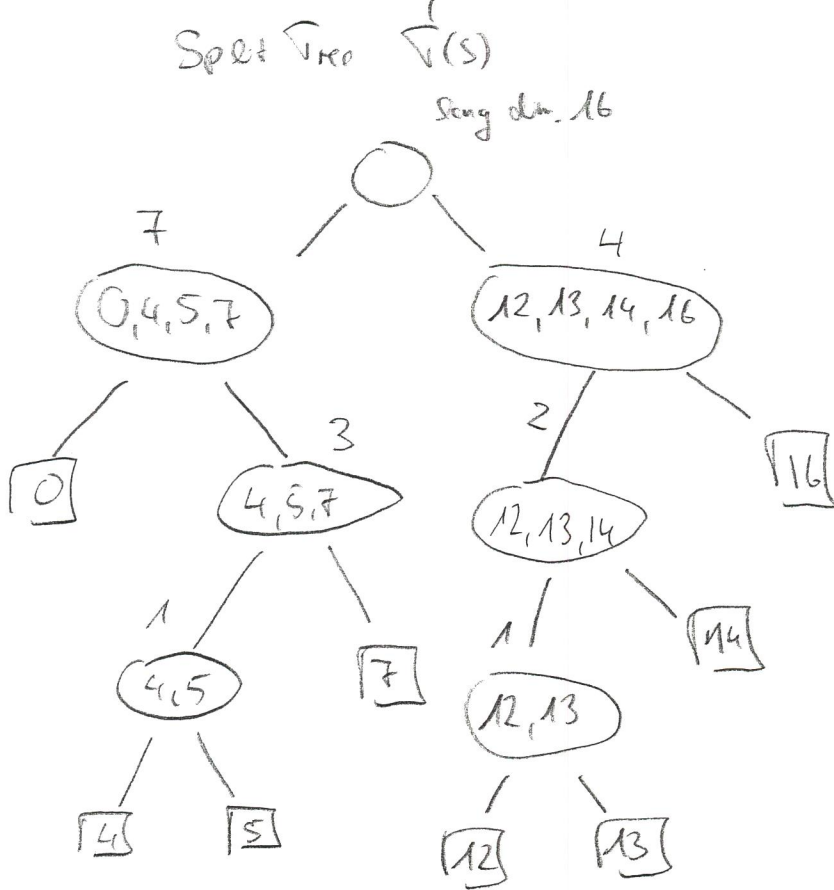
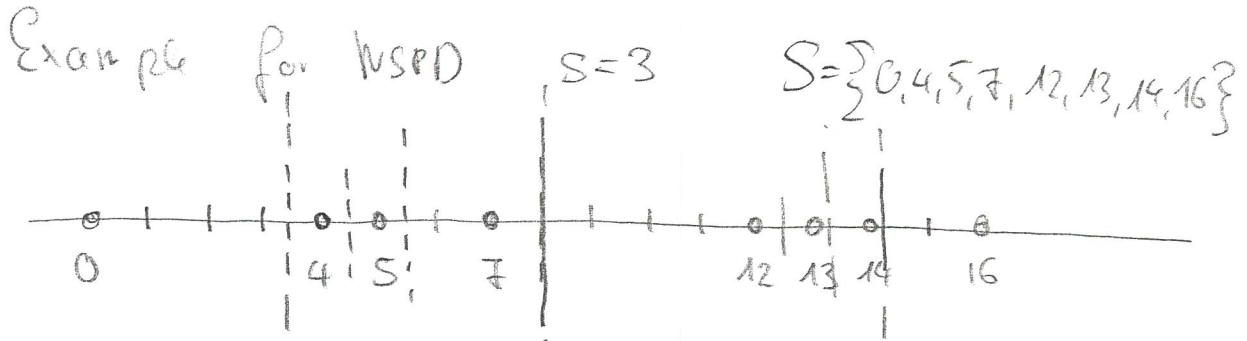
$$(a, b) (a, b') \Rightarrow S_b \cap S_{b'} = \emptyset$$

$$i, j \text{ leaves of } T \Rightarrow S_i \subset S_j \text{ or } S_j \subset S_i \text{ or } S_i \cap S_j = \emptyset$$

$\square$

Not clear: number of pairs,  $m$ , reported

There are only  $O(n)$  subsets  $S_a$  for the mins nodes  $a \in V$ , but the same set can occur many times as  $A_i, A_j, B_k, B_l$  etc.



$D_{min} = 1$

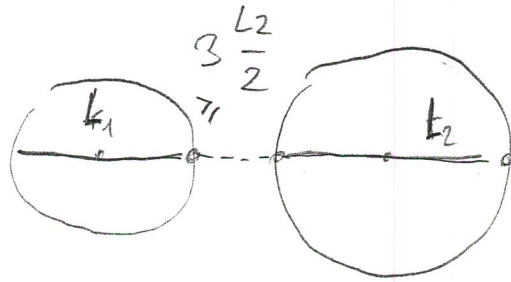
$L_{max} = \text{length of the set}$

$\frac{\sqrt{L}}{2}$

$\frac{L}{2}$  Radius of smallest enclosing "disk"

A, B well separated w.r.t  $s=3$

$\Leftrightarrow$



$L_1 < L_2$

Distance  $\geq \frac{3}{2}$  max length

WSPD - Computation via Find Pairs (Forest, Collection of Trees)

$FP(\{0, 4, 5, 7\}, \{12, 13, 14, 16\})$

$7 + 3 \cdot \frac{7}{2} > 12$

$FP(\{0\}, \{12, 13, 14, 16\})$   
 $0 + 3 \cdot \frac{4}{2} \leq 12$

$FP(\{4, 5, 7\}, \{12, 13, 14, 16\})$

$7 + 3 \cdot \frac{4}{2} > 5$

$FP(\{4, 5, 7\}, \{12, 13, 14\})$

$FP(\{4, 5, 7\}, \{16\})$

$7 + 3 \cdot \frac{3}{2} \leq 12$

$7 + 3 \cdot \frac{3}{2} \leq 12$

$FP(\{0\}, \{4, 5, 7\})$

$0 + 3 \cdot \frac{3}{2} > 4$

$FP(\{0\}, \{4, 5\})$  and  $FP(\{0\}, \{7\})$

$0 + 3 \cdot \frac{1}{2} \leq 4$

$FP(\{4, 5\}, \{7\})$

$5 + 3 \cdot \frac{1}{2} \leq 7$

$FP(\{12, 13\}, \{14\})$

$13 + 3 \cdot \frac{1}{2} > 14$

$FP(\{12\}, \{14\})$

$FP(\{13\}, \{14\})$

$FP(\{12, 13, 14\}, \{16\})$

$14 + 3 \cdot 1 > 16$

$FP(\{12, 13\}, \{16\})$

$13 + 3 \cdot \frac{1}{2} \leq 16$

$FP(\{14\}, \{16\})$

$FP(\{4\}, \{5\})$

$FP(\{12\}, \{13\})$